

COMP 520 - Compilers

Lecture 8 (Tue Feb 8, 2022)

Abstract Syntax Trees

- Reading for Feb 10
 - Finish Chapter 4
 - section 4.6 – Syntactic Analysis in Triangle (pp 124 – 128)
- Materials on our website
 - simpleAST
 - illustrates AST construction and evaluation for miniArith
 - AbstractSyntaxTrees.zip
 - `miniJava.AbstractSyntaxTrees` package needed to construct AST's for miniJava in PA2

Topics

- From concrete syntax trees to abstract syntax trees
 - AST “grammar”
 - AST representation choices
 - AST construction
- AST traversal
 - generalize using Visitor
- miniJava AST classes (available on web)
 - use these to construct the AST in PA2
 - you will need to provide some implementation of sourcePosition



Concrete syntax tree

- Concrete syntax is described by an (EBNF) CFG grammar

– ex: simple arithmetic expressions

$S ::= E \$$

$E ::= E \text{ Op } T \mid T$

$T ::= (E) \mid \text{Num}$

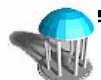
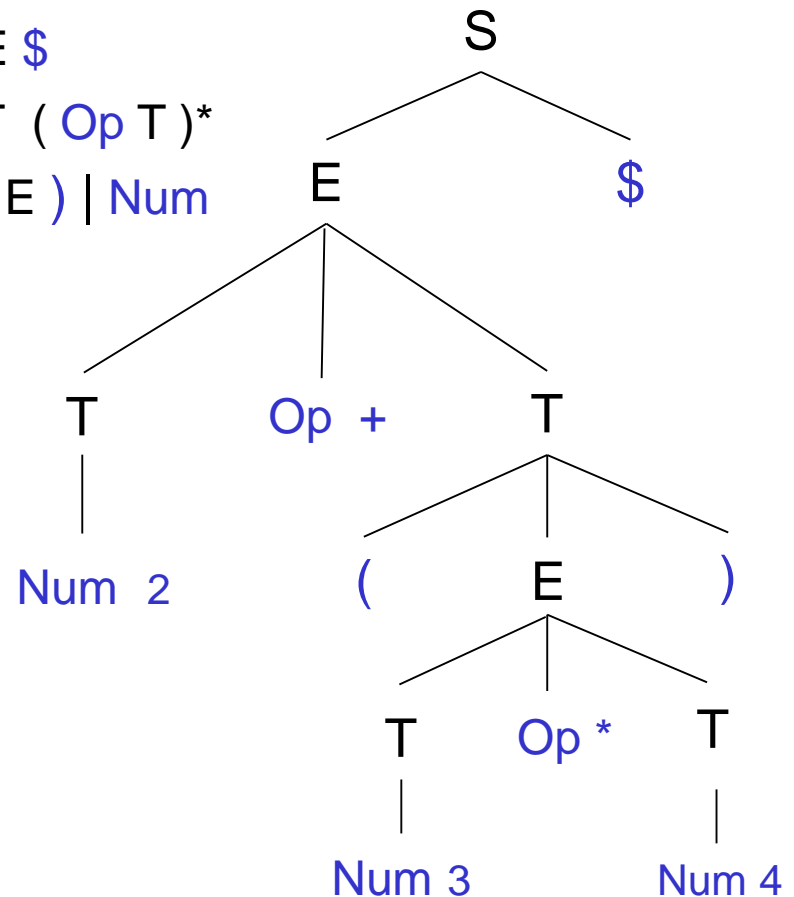


$S ::= E \$$

$E ::= T (\text{Op } T)^*$

$T ::= (E) \mid \text{Num}$

- Concrete syntax tree for
 $2 + (3 * 4) \$$

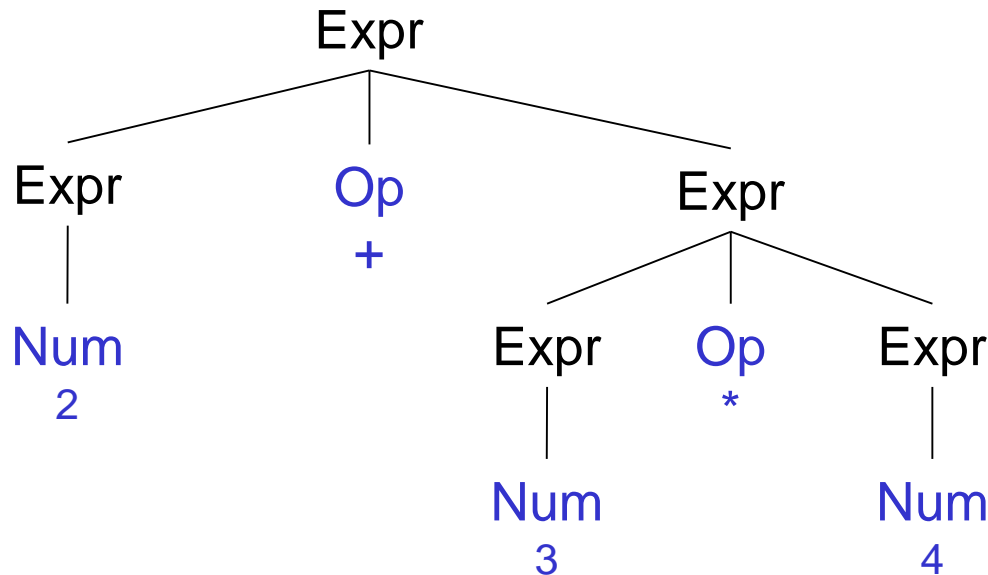


Abstract syntax

- Abstract syntax can be described by a simple CFG grammar
 - ex: simple arithmetic expressions

Expr ::= Expr Op Expr (BinExpr)
| Num (NumExpr)

- Abstract syntax tree for $2 + (3 * 4)$



AST representation in Java (Strategy 1)

- AST classes define an (n-ary) tree with tagged nodes
 - simple to define, more complex to construct and traverse
 - minimal benefit from Java type checker

```
enum NodeType { UNARYEXPR, BINARYEXPR, NUM }
public abstract class AST {
    public NodeType n;
}

public class Node extends AST {
    public AST[] children;
    public Node(NodeType n, AST[] children) { ... }
}

public class Leaf extends AST {
    public String spelling;
    public Leaf(NodeType n, String spelling) { ... }
}
```



AST representation in Java (Strategy 2)

- **AST classes closely follow structure of AST grammar**
 - more complex to define AST classes, but simpler to construct AST
 - we gain considerable benefit from the Java type checker
 - more rigorously supports AST traversal using a Visitor interface
- **Rules**
 - create abstract class AST
 - for each nonterminal in AST grammar
 - construct an abstract subclass of AST
 - for each choice within a rule
 - construct a concrete subclass of the LHS nonterminal
- We will follow this strategy



AST representation in Java (Strategy 2)

- Example

Expr ::= Expr Oper Expr (BinExpr)
| Num (NumExpr)

```
abstract public class AST {}

abstract public class Expr extends AST {}

public class BinExpr extends Expr {
    public Terminal oper;
    public Expr left, right;
    public BinExpr(Expr left, Terminal oper, Expr right) { ... }
}

public class NumExpr extends Expr {
    public Terminal num;
    public NumExpr(Terminal num) { ... }
}

public class Terminal extends AST {
    public String spelling;
    public Terminal(String spelling) { ... };
}
```



Building an AST during a concrete syntax parse

- concrete syntax for arithmetic expression grammar

$$E ::= T \mid E \textit{ op} T$$
$$T ::= (E) \mid \textit{ num}$$

- transformed and augmented

$$S ::= E \$$$
$$E ::= T (\textit{ op} T)^*$$
$$T ::= (E) \mid \textit{ num}$$

- abstract syntax

$$\begin{array}{l} \text{Expr} ::= \text{Expr Op Expr} \quad (\text{BinExpr}) \\ \quad \quad \quad \mid \text{Num} \quad \quad \quad (\text{NumExpr}) \end{array}$$

- how to build AST?

- modify parse procedures to return pieces of AST

```
Expr parseS() {
  Expr e = parseE();
  accept(Token.eot);
  return e
}

Expr parseE() {
  Expr e1 = parseT();
  while (curToken.kind == Token.op) {
    Terminal op = new Terminal(curToken);
    acceptIt();
    Expr e2 = parseT();
    e1 = new BinExpr(e1,op,e2);
  }
  return e1;
}

Expr parseT() {
  case (curToken.kind) {
    Token.LPAREN:
      acceptIt();
      Expr e1 = parseE();
      accept(Token.RPAREN);
      return e1;

    Token.num:
      NumExpr e2 = new NumExpr(curToken);
      acceptIt();
      return e2;
  }
}
```



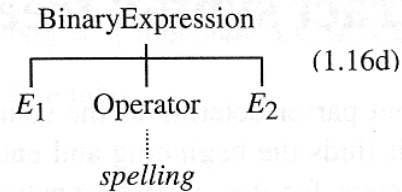
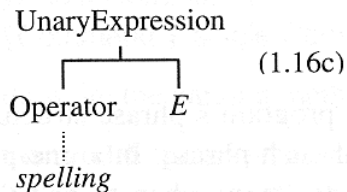
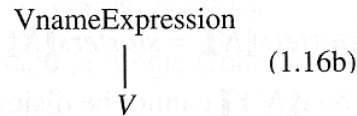
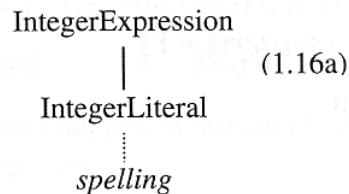
mini-Triangle – Expression ASTs

Concrete Syntax (EBNF)

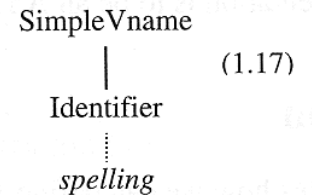
Tokens

Expression	::=	primary-Expression	(1.4a)
		Expression Operator primary-Expression	(1.4b)
primary-Expression	::=	Integer-Literal	(1.5a)
		V-name	(1.5b)
		Operator primary-Expression	(1.5c)
		(Expression)	(1.5d)
V-name	::=	Identifier	(1.6)
Operator	::=	+ - * / < > = \	(1.10a–h)
Identifier	::=	Letter Identifier Letter Identifier Digit	(1.11a–c)
Integer-Literal	::=	Digit Integer-Literal Digit	(1.12a–b)

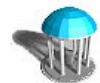
• Expression ASTs (E):



• V-name ASTs (V):



Abstract Syntax (ASTs)



mini-Triangle – Command ASTs

Concrete Syntax

Program	::=	single-Command	(1.1)
Command	::=	single-Command	(1.2a)
		Command ; single-Command	(1.2b)
single-Command	::=	V-name := Expression	(1.3a)
		Identifier (Expression)	(1.3b)
		if Expression then single-Command	(1.3c)
		else single-Command	
		while Expression do single-Command	(1.3d)
		let Declaration in single-Command	(1.3e)
begin Command end	(1.3f)		

ASTs



mini-Triangle – Declaration ASTs

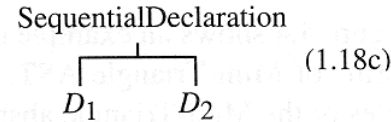
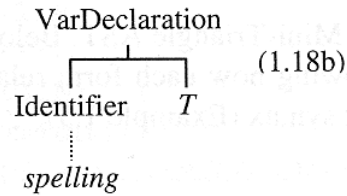
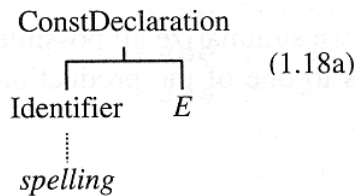
Concrete Syntax

Declaration ::= single-Declaration (1.7a)
 | Declaration ; single-Declaration (1.7b)

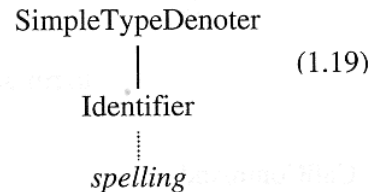
single-Declaration ::= **const** Identifier ~ Expression (1.8a)
 | **var** Identifier : Type-denoter (1.8b)

Type-denoter ::= Identifier (1.9)

- Declaration ASTs (D):



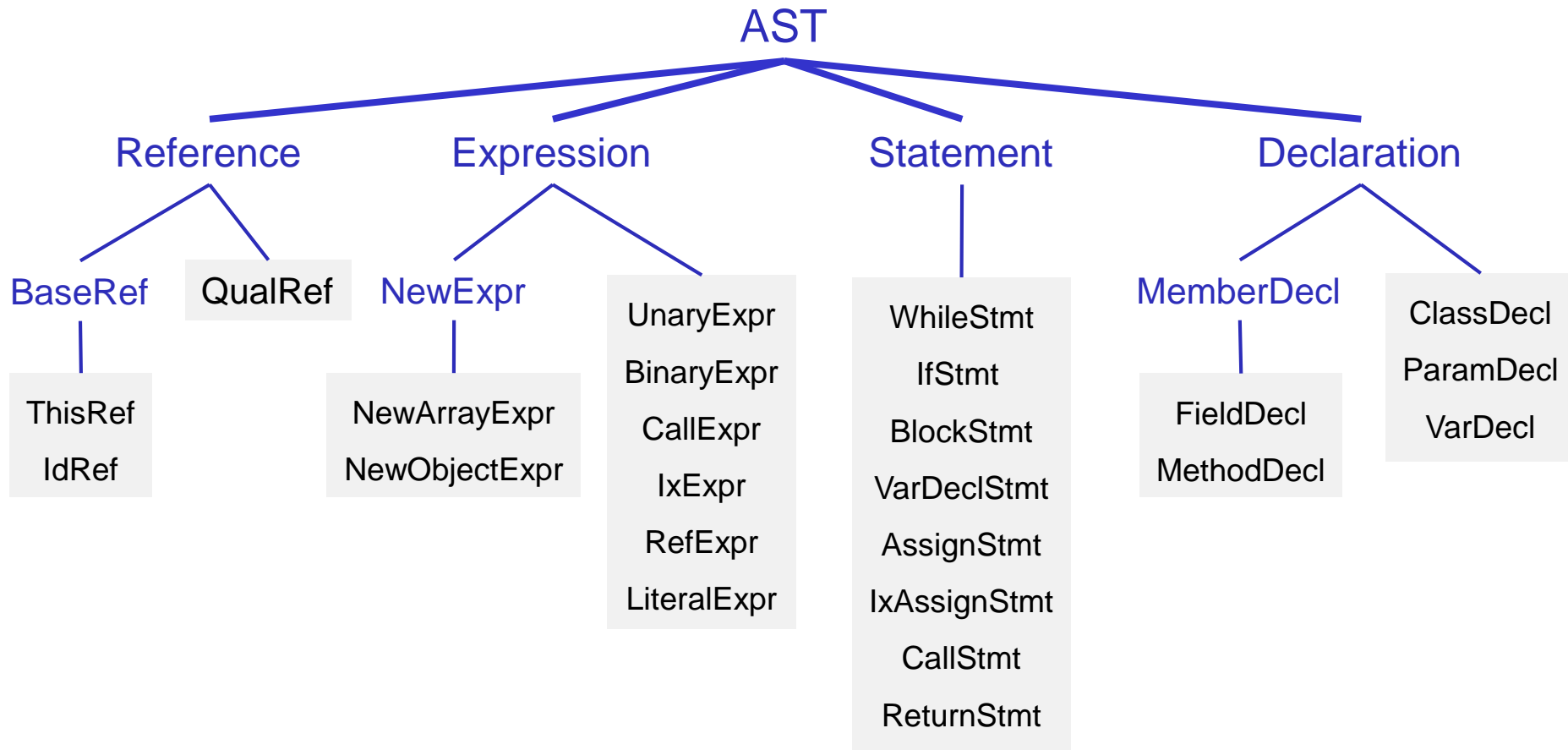
- Type-denoter ASTs (T):



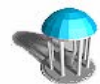
ASTs



The miniJava AST classes



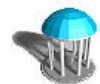
Additional constructors for lists of class instances, with an iterator for traversal of the list:
ClassDeclList, FieldDeclList, MethodDeclList, ParamDeclList, StatementList, ExprList





miniJava AST types

- ✓ **ⓐ** AST
 - ✓ **ⓐ** Declaration
 - ⓐ ClassDecl
 - ✓ **ⓐ** LocalDecl
 - ⓐ ParameterDecl
 - ⓐ VarDecl
 - ✓ **ⓐ** MemberDecl
 - ⓐ FieldDecl
 - ⓐ MethodDecl
 - ✓ **ⓐ** Expression
 - ⓐ BinaryExpr
 - ⓐ CallExpr
 - ⓐ lxEpr
 - ⓐ LiteralExpr
 - ✓ **ⓐ** NewExpr
 - ⓐ NewArrayExpr
 - ⓐ NewObjectExpr
 - ⓐ RefExpr
 - ⓐ UnaryExpr
 - ⓐ Package
- ✓ **ⓐ** Reference
 - ✓ **ⓐ** BaseRef
 - ⓐ IdRef
 - ⓐ ThisRef
 - ⓐ QualRef
- ✓ **ⓐ** Statement
 - ⓐ AssignStmt
 - ⓐ BlockStmt
 - ⓐ CallStmt
 - ⓐ IfStmt
 - ⓐ lxEprAssignStmt
 - ⓐ ReturnStmt
 - ⓐ VarDeclStmt
 - ⓐ WhileStmt
- ✓ **ⓐ** Terminal
 - ⓐ BooleanLiteral
 - ⓐ Identifier
 - ⓐ IntLiteral
 - ⓐ Operator
- ✓ **ⓐ** TypeDenoter
 - ⓐ ArrayType
 - ⓐ BaseType
 - ⓐ ClassType
- ⓐ ASTDisplay



Implementing a general AST traversal

- **Strategy 1**
 - add methods to each AST class for each kind of traversal
 - **Example**
 - display methods for AST display
 - eval methods for AST evaluation
 - drawback
 - **Classes become very large when traversals get complex**
 - Contextual analysis
 - Code generation
 - **Code for each kind of traversal is scattered over many classes**



Implementing a general AST traversal

- **Strategy 2**
 - use a *visitor* pattern
 - A **visitor Interface**
 - specifies a visitX method for each concrete class X in AST
 - a **visit method** in each AST class
 - public Object visit(Visitor v, Object arg)
 - each separate traversal implements the Visitor interface
 - **sample Visitor instances**
 - AST display
 - Identification
 - Type checking
 - Code generation
 - **Code for each type of visitor is collected in one class**
 - but a bit cumbersome to write out



Simple set of AST classes

```
abstract class AST {}
```

```
abstract class Expr extends AST {}
```

```
class BinExpr extends Expr {
```

```
    public Token oper;
```

```
    public Expr left;
```

```
    public Expr right;
```

```
    public BinExpr(Expr left, Token oper, Expr right) { . . . }  
}
```

```
class NumExpr extends Expr {
```

```
    public int val;
```

```
    public NumExpr(int val) { . . . }  
}
```



The Visitor interface

- Visitor interface requires a **visitX** method for every (non-abstract) AST class X

```
public interface Visitor {  
    visitBinaryExpr(BinaryExpr e);  
    visitNumExpr(NumExpr e);  
}
```

- Each AST class is augmented with a single **visit** method

```
class NumExpr extends Expr {  
    public int val;  
    public NumExpr(int v) { . . . }  
  
    public void visit(Visitor v) { v.visitNumExpr(this); }  
}
```

- All AST traversals use the same “**visit**” method in each node type
 - the visit method “connects” a specific visitor v to **this** specific node



Inorder traversal of the AST

- The **InorderWalk** AST traversal implements the **Visitor** interface

```
public class InorderWalk implements Visitor {
    public void visitBinExpr(BinExpr be) {
        be.left.visit(this);
        System.out.println(be.oper.spelling);
        be.right.visit(this);
    }

    public void visitNumExpr(NumExpr ne) {
        System.out.println(ne.val);
    }

    // print nodes of AST inorder
    public void walk(AST a) {
        a.visit(this);
    }
}
```



Adding arguments to the traversal (Book method)

- methods implemented by a Visitor have an **Object** arg and **Object** result

```
public interface Visitor {  
    Object visitBinExpr(BinExpr e, Object arg);  
    Object visitNumExpr(NumExpr e, Object arg);  
}
```

- AST class **visit** method has an **Object** arg and **Object** result

```
class NumExpr extends Expr {  
    public int val;  
  
    public NumExpr(int val) { . . . }  
  
    public Object visit(Visitor v, Object arg) {  
        return v.visitNumExpr(this, arg);  
    }  
}
```



Example use of the Visitor

- Individual traversals implement Visitor with custom logic for each node type

```
class Checker implements Visitor {  
  
    public Object visitAssignStmt(AssignStmt s, Object arg) {  
        Type t1 = (Type) s.ref.visit(this, arg);  
        Type t2 = (Type) s.exp.visit(this, arg);  
        return (t1.equals(t2)? t1 : Type.ERROR)  
    }  
}
```

- Good solution?
 - (+) appropriately OO
 - (+) compiler insures visitor defined for all node types
 - (+) specific definitions gathered together in a single class
 - (-) Object parameters and results will require a lot of explicit casting



Adding arguments to the traversal (parameterized types)

```
public interface Visitor<ArgType, ResType> {  
    public ResType visitBinExpr(BinExpr expr, ArgType arg);  
    public ResType visitNumExpr(NumExpr expr, ArgType arg);  
}
```

```
class NumExpr extends Expr {  
    public int val;  
    public NumExpr(int val) { . . . }  
  
    public <ArgType, ResType> ResType  
        visit(Visitor<ArgType, ResType> v, ArgType arg) {  
            return v.visitNumExpr(this, arg);  
        }  
}
```



Example use of the Visitor

- Individual traversals implement Visitor with custom logic for each node type

```
class Checker implements Visitor<Type, Type> {  
  
    ...  
  
    public Type visitAssignStmt(AssignStmt s, Type arg) {  
        Type t1 = s.ref.visit(this, arg);  
        Type t2 = s.exp.visit(this, arg);  
        if (! t1.equals(t2))  
            typeError("incompatible types in assignment", s);  
        return Types.StmtType;  
    }  
}
```

- Good solution?
 - Improves type checking in the visitors
 - Improves readability



PA2 Submission Details

1. Your compiler project must follow these requirements with respect to the package names and structure
2. Your `Compiler.java` mainclass must have the functionality called out on the following page, but you may vary your implementation.




```

package miniJava;
import java.io.FileInputStream;
import java.io.FileNotFoundException;
import java.io.InputStream;

import miniJava.SyntacticAnalyzer.Parser;
import miniJava.SyntacticAnalyzer.Scanner;
import miniJava.AbstractSyntaxTrees.*;

/**
 * Recognize whether input file named in args[0] contains a syntactically valid
 * miniJava program, and, if valid, display corresponding AST.
 */
public class Compiler {

    public static void main(String[] args) {
        InputStream inputStream = null;
        try {
            inputStream = new FileInputStream(args[0]);
        }
        catch (FileNotFoundException e) {
            System.out.println("Input file " + args[0] + " not found");
            System.exit(3);
        }

        ErrorReporter errorReporter = new ErrorReporter();
        Scanner scanner = new Scanner(inputStream, errorReporter);
        Parser parser = new Parser(scanner, errorReporter);

        System.out.println("Syntactic analysis ... ");
        AST ast = parser.parseProgram();
        System.out.print("Syntactic analysis complete: ");
        if (errorReporter.hasErrors()) {
            System.out.println("Invalid miniJava program");
            System.exit(4);
        }
        else {
            System.out.println("Valid miniJava program:");
            new ASTDisplay().showTree(ast);
            System.exit(0);
        }
    }
}

```

use this termination code if unable to open input file

use these termination codes for invalid/valid miniJava programs, respectively

create AST

output AST using ASTDisplay(), for accepted miniJava programs

